

# Interpolation \*

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## 1 Approximation Theory

Given  $f \in C[a, b]$ , we would like to find a “good” approximation to it by “simpler functions”, i.e. functions in a given class (or family)  $\Phi$ . For example,  $\Phi = \mathbb{P}_n = \{\text{all polynomials of degree } \leq n\}$ .

A natural problem is that of finding the *best approximation* to  $f$  by functions in  $\Phi$ . But how do we measure the accuracy of any approximation? that is, what norm<sup>1</sup> do we use? We have several choices for norms of functions. The most commonly used are:

1. The max or infinity norm:  $\|f\|_\infty = \sup_{x \in [a, b]} |f(x)|$ .

2. The 2-norm:  $\|f\|_2 = \left(\int_a^b f^2(x) dx\right)^{\frac{1}{2}}$ .

3. The p-norm:  $\|f\|_p = \left(\int_a^b f^p(x) dx\right)^{\frac{1}{p}}$ .

Later, we will need to consider weighted norms: for some positive function  $\omega(x)$  in  $[a, b]$  (it could be zero on a finite number of points) we define

$$(1) \quad \|f\|_{\omega, 2} = \left(\int_a^b \omega(x) f^2(x) dx\right)^{\frac{1}{2}}.$$

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<sup>1</sup>A norm  $\|\cdot\|$  is a real valued function on a vector space such that (1)  $\|f\| > 0, f \neq 0$ , (2)  $\|\lambda f\| = |\lambda| \|f\|$  for any  $\lambda$  scalar, and (3)  $\|f + g\| \leq \|f\| + \|g\|$ .

Then, by best approximation in  $\Phi$  we mean a function  $p \in \Phi$  such that

$$\|f - p\| \leq \|f - q\|, \forall q \in \Phi.$$

Computationally, it is often more efficient to seek not the best approximation but one that is sufficiently accurate and fast converging to  $f$ . The central building block for this approximation is the problem of interpolation.

## 2 Interpolation

Let us focus on the case of approximating a given function by a polynomial of degree at most  $n$ . Then the *interpolation problem* can be stated as follows: Given  $n+1$  distinct points,  $x_0, x_1, \dots, x_n$  called *nodes* and corresponding values  $f(x_0), f(x_1), \dots, f(x_n)$ , find a polynomial of degree at most  $n$ ,  $P_n(x)$ , which satisfies (the interpolation property)

$$\begin{aligned} P_n(x_0) &= f(x_0) \\ P_n(x_1) &= f(x_1) \\ &\vdots \\ P_n(x_n) &= f(x_n). \end{aligned}$$

Let us represent such polynomial as  $P_n(x) = a_0 + a_1x + \dots + a_nx^n$ . Then, the interpolation property means

$$P_n(x_0) = f(x_0), P_n(x_1) = f(x_1), \dots, P_n(x_n) = f(x_n),$$

which implies

$$\begin{aligned} a_0 + a_1x_0 + \dots + a_nx_0^n &= f(x_0) \\ a_0 + a_1x_1 + \dots + a_nx_1^n &= f(x_1) \\ &\vdots \\ a_0 + a_1x_n + \dots + a_nx_n^n &= f(x_n). \end{aligned}$$

This is a linear system of  $n+1$  equations in  $n+1$  unknowns (the polynomial coefficients  $a_0, a_1, \dots, a_n$ ). In matrix form:

$$(2) \quad \begin{bmatrix} 1 & x_0 & x_0^2 & \dots & x_0^n \\ 1 & x_1 & x_1^2 & \dots & x_1^n \\ \vdots & & & & \\ 1 & x_n & x_n^2 & \dots & x_n^n \end{bmatrix} \begin{bmatrix} a_0 \\ a_1 \\ \vdots \\ a_n \end{bmatrix} = \begin{bmatrix} f(x_0) \\ f(x_1) \\ \vdots \\ f(x_n) \end{bmatrix}$$

Does this linear system have a solution? Is this solution unique? The answer is yes to both. Here is a simple proof. Take  $f \equiv 0$ , then  $P_n(x_j) = 0$ , for  $j = 0, 1, \dots, n$  but  $P_n$  is a polynomial of degree  $\leq n$ , it cannot have  $n + 1$  zeros unless  $P_n(x) \equiv 0$ , which implies  $a_0 = a_1 = \dots = a_n = 0$ . That is, the homogenous problem associated with (2) has only the trivial solution. Therefore, (2) has a unique solution.

In general the values to interpolate might not come from a function. They are just data supplied to us. We will often write  $(x_0, f_0), (x_1, f_1)$ , etc., to emphasize this more general setting.

**Example 1.** *As an illustration let us consider interpolation by a linear polynomial,  $P_1(x)$ . Suppose we are given  $(x_0, f_0)$  and  $(x_1, f_1)$ . We have written  $P_1(x)$  explicitly in the Introduction. We write it now in a different form:*

$$(3) \quad P_1(x) = \frac{x - x_1}{x_0 - x_1} f_0 + \frac{x - x_0}{x_1 - x_0} f_1$$

*Clearly, this polynomial has degree at most 1 and satisfies the interpolation property:*

$$(4) \quad P_1(x_0) = f_0,$$

$$(5) \quad P_1(x_1) = f_1.$$

**Example 2.** *Given  $(x_0, f_0), (x_1, f_1), (x_2, f_2)$  let us construct  $P_2(x)$ , the polynomial of degree at most 2 which interpolates these points. The way we have written  $P_1(x)$  in (3) is suggestive of how to explicitly write  $P_2(x)$ :*

$$P_2(x) = \frac{(x - x_1)(x - x_2)}{(x_0 - x_1)(x_0 - x_2)} f_0 + \frac{(x - x_0)(x - x_2)}{(x_1 - x_0)(x_1 - x_2)} f_1 + \frac{(x - x_0)(x - x_1)}{(x_2 - x_0)(x_2 - x_1)} f_2.$$

*If we define*

$$(6) \quad l_0^{(2)}(x) = \frac{(x - x_1)(x - x_2)}{(x_0 - x_1)(x_0 - x_2)},$$

$$(7) \quad l_1^{(2)}(x) = \frac{(x - x_0)(x - x_2)}{(x_1 - x_0)(x_1 - x_2)},$$

$$(8) \quad l_2^{(2)}(x) = \frac{(x - x_0)(x - x_1)}{(x_2 - x_0)(x_2 - x_1)},$$

*then we simply have*

$$(9) \quad P_2(x) = l_0^{(2)}(x) f_0 + l_1^{(2)}(x) f_1 + l_2^{(2)}(x) f_2.$$

Note that each of the polynomials (6), (7), and (8) are exactly of degree 2 and they satisfy  $l_j^{(2)}(x_k) = \delta_{jk}$ <sup>2</sup>. Therefore, it follows that  $P_2(x)$  given by (9) satisfies the interpolation property

$$(10) \quad P_2(x_0) = f_0,$$

$$(11) \quad P_2(x_1) = f_1,$$

$$(12) \quad P_2(x_2) = f_2.$$

We can now write down the polynomial (of degree at most  $n$ ) which interpolates  $n + 1$  given values,  $(x_0, f_0), \dots, (x_n, f_n)$ , where the interpolation nodes  $x_0, \dots, x_n$  are assumed distinct.

Define

$$(13) \quad \begin{aligned} l_j^{(n)}(x) &= \frac{(x - x_0) \cdots (x - x_{j-1})(x - x_{j+1} \cdots (x - x_n))}{(x_j - x_0) \cdots (x_j - x_{j-1})(x_j - x_{j+1} \cdots (x_j - x_n))} \\ &= \prod_{k=0, k \neq j}^n \frac{(x - x_k)}{(x_j - x_k)}, \quad \text{for } j = 0, 1, \dots, n. \end{aligned}$$

These are called the *elementary Lagrange polynomials* of degree  $n$ . Note that  $l_j^{(n)}(x_k) = \delta_{jk}$ . Therefore

$$(14) \quad P_n(x) = l_0^{(n)}(x)f_0 + l_1^{(n)}(x)f_1 + \cdots + l_n^{(n)}(x)f_n = \sum_{j=0}^n l_j^{(n)}(x)f_j$$

interpolates the given data, i.e., it satisfies the interpolation property  $P_n(x_j) = f_j$  for  $j = 0, 1, 2, \dots, n$ . Relation (14) is called the *Lagrange form* of the interpolating polynomial. The following result summarizes our discussion.

**Theorem 1.** *Given the  $n + 1$  values  $(x_0, f_0), \dots, (x_n, f_n)$ , for  $x_0, x_1, \dots, x_n$  distinct. There is a unique polynomial of degree at most  $n$ ,  $P_n(x)$ , such that  $P_n(x_j) = f_j$  for  $j = 0, 1, \dots, n$ .*

*Proof.*  $P_n(x)$  in (14) is of degree at most  $n$  and interpolates the data. Uniqueness follows from the fundamental algebra : suppose there is another polynomial  $Q_n(x)$  of degree at most  $n$  such that  $Q_n(x_j) = f_j$  for  $j = 0, 1, \dots, n$ . Consider  $W(x) = P_n(x) - Q_n(x)$ . This is a polynomial of degree at most  $n$  and  $W(x_j) = P_n(x_j) - Q_n(x_j) = f_j - f_j = 0$  for  $j = 0, 1, 2, \dots, n$ , which is impossible unless  $W(x) \equiv 0$  which implies  $Q_n = P_n$ .  $\square$

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<sup>2</sup> $\delta_{jk}$  is the Kronecker delta, i.e.  $\delta_{jk} = 0$  if  $k \neq j$  and 1 if  $k = j$ .

### 3 Connection to Best Approximation

We can view interpolation as a linear operator. Suppose that we have  $n + 1$  distinct nodes  $x_0, x_1, \dots, x_n$  contained in an interval  $[a, b]$ . Let  $f$  and  $g$  be two continuous functions in  $[a, b]$  and  $\alpha$  and  $\beta$  two scalars. Then, the interpolating polynomial for  $\alpha f(x) + \beta g(x)$  is  $P(x) = \alpha P_n(x) + \beta Q_n(x)$  where  $P_n(x)$  and  $Q_n(x)$  are the interpolating polynomials of  $f$  and  $g$ , respectively. This follows immediately from (14). Also, note that if  $f$  is a polynomial of degree at most  $n$ , its interpolating polynomial is itself, i.e.  $P_n(x) = f(x)$ .

Now suppose that  $P_n^*(x)$  is the best polynomial approximation of  $f$  in the max or uniform norm, i.e.

$$(15) \quad \min_{p \in \mathbb{P}_n} \|f - p\|_\infty = \|f - P_n^*\|_\infty,$$

where  $\mathbb{P}_n = \{\text{all polynomials of degree} \leq n\}$ . Let  $P_n(x)$  be the interpolating polynomials of  $f$  at  $x_0, x_1, \dots, x_n$ . Then,

$$\|f - P_n\|_\infty = \|f - P_n^* - (P_n - P_n^*)\|_\infty \leq \|f - P_n^*\|_\infty + \|P_n - P_n^*\|_\infty.$$

But  $P_n(x) - P_n^*(x)$  is a polynomial of degree at most  $n$  which interpolates  $f - P_n^*$ , that is

$$P_n(x) - P_n^*(x) = \sum_{j=0}^n l_j^{(n)}(x)(f(x_j) - P_n^*(x_j)).$$

Therefore,

$$(16) \quad \|P_n - P_n^*\|_\infty \leq \Lambda_n \|f - P_n^*\|_\infty$$

where

$$(17) \quad \Lambda_n = \max_{a \leq x \leq b} \sum_{j=0}^n |l_j^{(n)}(x)|$$

is called the *Lebesgue Constant*. Using this in (16) we obtain

$$(18) \quad \|f - P_n\|_\infty \leq (1 + \Lambda_n) \|f - P_n^*\|_\infty.$$

This inequality connects the interpolation error  $\|f - P_n\|_\infty$  with the best approximation error  $\|f - P_n^*\|_\infty$ . Let us see if can extract more information from this connection.

There is a fundamental result in approximation theory, which states that any continuous function can be approximated uniformly, with arbitrary accuracy by a polynomial. This is the celebrated Weierstrass Theorem.

**Theorem 2.** *Weierstrass Theorem.* Let  $f$  be a continuous function in  $[a, b]$ . Given  $\epsilon > 0$  there is a polynomial  $P$  such that

$$\|f - P\|_\infty < \epsilon.$$

Weierstrass theorem implies that as we increase the degree the best approximation polynomial converges uniformly to  $f$ , that is  $\|f - P_n^*\|_\infty \rightarrow 0$  as  $n \rightarrow \infty$ . However, because the Lebesgue constant  $\Lambda_n$  is not bounded in  $n$ , we cannot conclude that  $\|f - P_n\|_\infty \rightarrow 0$  as  $n \rightarrow \infty$ , i.e. that the interpolating polynomial, as we add more and more nodes, converges uniformly to  $f$ . That depends on the regularity of  $f$  and on the distribution of the nodes. We will discuss this further later.

## 4 Barycentric Formula

The Lagrange form of the interpolating polynomial is not convenient for computations. If we want to increase the degree of the polynomial we cannot reuse the work done in getting and evaluating a lower degree one. However, we can obtain a very efficient formula by rewriting the interpolating polynomial as follows.

Let

$$(19) \quad w(x) = (x - x_0)(x - x_1) \cdots (x - x_n).$$

The numerator in  $l_j^{(n)}(x)$  is  $w(x)/(x - x_j)$ . Thus, we can write

$$(20) \quad l_j^{(n)}(x) = w(x) \frac{\lambda_j^{(n)}}{x - x_j}, \quad \lambda_j^{(n)} = \frac{1}{\prod_{\substack{k=0 \\ k \neq j}}^n (x_j - x_k)}.$$

Therefore

$$(21) \quad P_n(x) = \sum_{j=0}^n w(x) \frac{\lambda_j^{(n)}}{x - x_j} f(x_j) = w(x) \sum_{j=0}^n \frac{\lambda_j^{(n)}}{x - x_j} f(x_j).$$

Now, note that from (14) with  $f(x) \equiv 1$  it follows that

$$(22) \quad \sum_{j=0}^n l_j^{(n)}(x) = 1$$

and diving (21) by (22) we get the so-called *Barycentric Formula* for interpolation:

$$(23) \quad P_n(x) = \frac{\sum_{j=0}^n \frac{\lambda_j^{(n)}}{x - x_j} f(x_j)}{\sum_{j=0}^n \frac{\lambda_j^{(n)}}{x - x_j}}, \quad \text{for } x \neq x_j, j = 0, 1, \dots, n.$$

For  $x = x_j$ ,  $j = 0, 1, \dots, n$ , the interpolation property should be used:  $P_n(x_j) = f(x_j)$ .

The numbers  $\lambda_j^{(n)}$  depend only on the nodes  $x_0, x_1, \dots, x_n$  and not on given values  $f(x_0), f(x_1), \dots, f(x_n)$ . We can precompute them efficiently as follows:

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 $\lambda_0^{(0)} = 1;$ 
for  $m = 1 : n$ 
  for  $j = 0 : m - 1$ 
     $\lambda_j^{(m)} = \frac{\lambda_j^{(m-1)}}{x_j - x_m};$ 
  end
   $\lambda_m^{(m)} = \frac{1}{\prod_{k=0}^{m-1} (x_m - x_k)};$ 
end

```

If we want to add one more point  $(x_{n+1}, f(x_{n+1}))$  we just extend the  $m$ -loop to  $n + 1$  to generate  $\lambda_0^{(n+1)}, \lambda_1^{(n+1)}, \dots, \lambda_{n+1}^{(n+1)}$ .

For equidistributed points,  $x_j = x_0 + jh$ ,  $j = 0, 1, \dots, n$  we have

$$\begin{aligned}
\lambda_j^{(n)} &= \frac{1}{(x_j - x_0) \cdots (x_j - x_{j-1})(x_j - x_{j+1}) \cdots (x_j - x_n)} \\
&= \frac{1}{(jh)[(j-1)h] \cdots (h)(-h)(-2h) \cdots (j-n)h} \\
&= \frac{1}{(-1)^{n-j} h^n [j(j-1) \cdots 1][1 \cdot 2 \cdots (n-j)]} \\
&= \frac{(-1)^{j-n}}{h^n n!} \frac{n!}{j!(n-j)!} \\
&= \frac{(-1)^{j-n}}{h^n n!} \binom{n}{j} \\
&= \frac{(-1)^n}{h^n n!} (-1)^j \binom{n}{j}.
\end{aligned}$$

We can omit the factor  $\frac{(-1)^n}{h^n n!}$  because it cancels out in the Barycentric Formula. Thus, for equidistributed points we can use

$$(24) \quad \lambda_j^{(n)} = (-1)^j \binom{n}{j}, \quad j = 0, 1, \dots, n.$$

## 5 Newton's Form and Divided Differences

There is another representation of the interpolating polynomial which is both very efficient computationally and very convenient in the derivation of numerical methods based on interpolation. The idea of this representation, due to Newton, is to use successively lower order polynomials for constructing  $P_n(x)$ .

Suppose we have gotten  $P_{n-1}(x)$ , that is the interpolating poly of  $f$  at  $x_0, \dots, x_{n-1}$  and we would like to obtain  $P_n(x)$ , the interpolating poly of  $f$  at  $x_0, \dots, x_n$  by reusing  $P_{n-1}(x)$ .  $P_n(x) - P_{n-1}(x) = R(x)$ , where  $R(x)$  is a polynomial of degree at most  $n$ . Moreover, for  $j = 0, \dots, n-1$

$$(25) \quad r(x_j) = P_n(x_j) - P_{n-1}(x_j) = f(x_j) - f(x_j) = 0.$$

Therefore,  $R(x)$  can be factored as

$$(26) \quad R(x) = c_n(x - x_0)(x - x_1) \cdots (x - x_{n-1}).$$



The constant  $c_n$  is called the  $n$ -th *divided difference* of  $f$  with respect to  $x_0, x_1, \dots, x_n$ , and is usually denoted as  $f[x_0, \dots, x_n]$ . Thus, we have

$$(27) \quad P_n(x) = P_{n-1}(x) + f[x_0, \dots, x_n](x - x_0)(x - x_1) \cdots (x - x_{n-1})$$

By the same argument, we have

$$(28) \quad P_{n-1}(x) = P_{n-2}(x) + f[x_0, \dots, x_{n-1}](x - x_0)(x - x_1) \cdots (x - x_{n-2}),$$

etc. So we arrive at Newton's Form of  $P_n(x)$ :

$$(29) \quad P_n(x) = f[x_0] + f[x_0, x_1](x - x_0) + \dots + f[x_0, \dots, x_n](x - x_0) \cdots (x - x_{n-1}).$$

Note that for  $n = 1$

$$\begin{aligned} P_1(x) &= f[x_0] + f[x_0, x_1](x - x_0) \\ P_1(x_0) &= f[x_0] = f(x_0) \\ P_1(x_1) &= f[x_0] + f[x_0, x_1](x_1 - x_0) = f(x_1) \end{aligned}$$

Therefore

$$(30) \quad f[x_0] = f(x_0)$$

$$(31) \quad f[x_0, x_1] = \frac{f(x_1) - f(x_0)}{x_1 - x_0}.$$

and

$$(32) \quad P_1(x) = f(x_0) + \frac{f(x_1) - f(x_0)}{x_1 - x_0}(x - x_0).$$

Define  $f[x_j] = f(x_j)$  for  $j = 0, 1, \dots, n$ . The following identity will allow us to compute all the required divided differences.

**Theorem 3.**

$$(33) \quad f[x_0, x_1, \dots, x_k] = \frac{f[x_1, x_2, \dots, x_k] - f[x_0, x_1, \dots, x_{k-1}]}{x_k - x_0}.$$

*Proof.* Let  $P_{k-1}$  be the interpolating polynomial of degree at most  $k-1$  of  $f$  at  $x_1, \dots, x_k$  and  $Q_{k-1}$  the interpolating polynomial of degree at most  $k-1$  of  $f$  at  $x_0, \dots, x_{k-1}$ . Then

$$(34) \quad P(x) = P_{k-1}(x) + \frac{x - x_k}{x_k - x_0} [P_{k-1}(x) - Q_{k-1}(x)].$$

is a polynomial of degree at most  $k$  and for  $j = 1, 2, \dots, k-1$

$$P_k(x_j) = f(x_j) + \frac{x_j - x_k}{x_k - x_0} [f(x_j) - f(x_0)] = f(x_j).$$

Moreover,  $P_k(x_0) = Q_{k-1}(x_0) = f(x_0)$  and  $P_k(x_k) = P_{k-1}(x_k) = f(x_k)$ . Therefore,  $P$  is the interpolation poly of  $f$  at  $x_0, \dots, x_k$ . The leading order coefficient of  $P_k(x)$  is  $f[x_0, \dots, x_k]$  and equating this with the leading order coefficient of  $P$ ,  $\frac{f[x_1, \dots, x_k] - f[x_0, x_1, \dots, x_{k-1}]}{x_k - x_0}$ , gives (33).  $\square$

To use (33) to obtain the divided difference we construct a table. Setting  $f_j = f(x_j)$  this divided difference table is built column by column as illustrated below for  $n = 3$ .

| $x_j$ | 0th order | 1th order     | 2th order          | 3th order               |
|-------|-----------|---------------|--------------------|-------------------------|
| $x_0$ | $f_0$     |               |                    |                         |
|       |           | $f[x_0, x_1]$ |                    |                         |
| $x_1$ | $f_1$     |               | $f[x_0, x_1, x_2]$ |                         |
|       |           | $f[x_1, x_2]$ |                    | $f[x_0, x_1, x_2, x_3]$ |
| $x_2$ | $f_2$     |               | $f[x_1, x_2, x_3]$ |                         |
|       |           | $f[x_2, x_3]$ |                    |                         |
| $x_3$ | $f_3$     |               |                    |                         |

**Example 3.** Let  $f(x) = 1 + x^2$  and  $x_j = j - 1$  for  $j = 0, \dots, 3$ . Then

| $x_j$ | 0th order         | 1th order                           | 2th order                           | 3th order         |
|-------|-------------------|-------------------------------------|-------------------------------------|-------------------|
| 0     | $\textcircled{1}$ |                                     |                                     |                   |
|       |                   | $\frac{2-1}{1-0} = \textcircled{1}$ |                                     |                   |
| 1     | 2                 |                                     | $\frac{3-1}{2-0} = \textcircled{1}$ |                   |
|       |                   | $\frac{6-2}{2-1} = 3$               |                                     | $\textcircled{0}$ |
| 2     | 5                 |                                     | $\frac{5-3}{3-1} = 1$               |                   |
|       |                   | $\frac{10-5}{3-2} = 5$              |                                     |                   |
| 3     | 10                |                                     |                                     |                   |

so  $P_3(x) = 1 + 1(x-0) + 1(x-0)(x-1) + 0(x-0)(x-1)(x-2) = 1 + x^2$ .

After computing the divided differences, we need to evaluate  $P_n$  at a given point  $x$ . This can be done efficiently by suitably factoring it. For example, for  $n = 3$  we have

$$\begin{aligned} P_3(x) &= c_0 + c_1(x - x_0) + c_2(x - x_0)(x - x_1) + c_3(x - x_0)(x - x_1)(x - x_2) \\ &= c_0 + (x - x_0) \{c_1 + (x - x_1)[c_2 + (x - x_2)c_3]\} \end{aligned}$$

For general  $n$  we can use the following *Horner*-like scheme to get  $p = P_n(x)$ :

```
p = c_n;
for k = n - 1 : 0
    p = c_k + (x - x_k) * p;
end
Note that
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$$(35) \quad P_k(x) = P_{k-1}(x) + f[x_0, \dots, x_k](x - x_0) \cdots (x - x_{k-1})$$

where  $P_{k-1}(x)$  is a polynomial of degree  $\leq k - 1$ . Then, the leading order coefficient of  $P_k(x)$  is  $f[x_0, \dots, x_k]$  and consequently  $P_k^{(k)}(x) = k!f[x_0, \dots, x_k]$  and thus

$$(36) \quad f[x_0, \dots, x_k] = \frac{1}{k!} P_k^{(k)}(x).$$

## 6 Cauchy's Remainder

In the Introduction we proved that if  $x_0, x_1$ , and  $x$  are in  $[a, b]$  then

$$f(x) - P_1(x) = \frac{1}{2} f''(\xi(x))(x - x_0)(x - x_1),$$

where  $\xi(x) \in (a, b)$ . The general result about the interpolation error is the following theorem:

**Theorem 4.** *Let  $f \in C^{n+1}[a, b]$ ,  $x_0, x_1, \dots, x_n, x$  be contained in  $[a, b]$ , and  $P_n(x)$  be the interpolation polynomial of degree  $\leq n$  of  $f$  at  $x_0, \dots, x_n$  then*

$$(37) \quad f(x) - P_n(x) = \frac{1}{(n+1)!} f^{(n+1)}(\xi(x))(x - x_0)(x - x_1) \cdots (x - x_n),$$

where  $\min\{x_0, \dots, x_n, x\} < \xi(x) < \max\{x_0, \dots, x_n, x\}$ .

*Proof.* The right hand side of (37) is known as the Cauchy Remainder and the following proof is due to Cauchy.

For  $x$  equal to one of the nodes  $x_j$  the result is trivially true. Take  $x$  fixed not equal to any of the nodes and define

$$(38) \quad \phi(t) = f(t) - P_n(t) - [f(x) - P_n(x)] \frac{(t - x_0)(t - x_1) \cdots (t - x_n)}{(x - x_0)(x - x_1) \cdots (x - x_n)}.$$

Clearly,  $\phi \in C^{n+1}[a, b]$  and vanishes at  $t = x_0, x_1, \dots, x_n, x$ . That is,  $\phi$  has at least  $n + 2$  zeros. Applying Rolle's Theorem  $n + 1$  times we conclude that there exists a point  $\xi(x) \in (a, b)$  such that  $\phi^{(n+1)}(\xi(x)) = 0$ . Therefore,

$$0 = \phi^{(n+1)}(\xi(x)) = f^{(n+1)}(\xi(x)) - [f(x) - P_n(x)] \frac{(n+1)!}{(x - x_0)(x - x_1) \cdots (x - x_n)}$$

from which (37) follows. Note that the repeated application of Rolle's theorem implies that  $\xi(x)$  is between  $\min\{x_0, x_1, \dots, x_n, x\}$  and  $\max\{x_0, x_1, \dots, x_n, x\}$ .  $\square$

We can use the same argument to relate divided differences to the derivatives of  $f$ . Let  $P_{n+1}(t)$  be the interpolating polynomial of  $f$  at  $x_0, \dots, x_n, x$ . Then  $f(t) - P_{n+1}(t)$  vanishes at  $t = x_0, \dots, x_n, x$ . By repeated application of Rolle's theorem there is  $\xi(x)$  between  $\min\{x_0, x_1, \dots, x_n, x\}$  and  $\max\{x_0, x_1, \dots, x_n, x\}$  such that

$$0 = f^{n+1}(\xi(x)) - P_{n+1}^{n+1}(\xi) = f^{n+1}(\xi(x)) - (n+1)!f[x_0, \dots, x_n, x].$$

Therefore

$$(39) \quad f[x_0, \dots, x_n, x] = \frac{f^{(n+1)}(\xi(x))}{(n+1)!}.$$

Similarly, if  $P_k$  is the interpolating polynomial of  $f$  at  $x_0, x_1, \dots, x_k$ , then

$$(40) \quad f[x_0, \dots, x_k] = \frac{1}{k!} f^{(k)}(\xi),$$

where  $\min\{x_0, \dots, x_k\} < \xi < \max\{x_0, \dots, x_k\}$ . Suppose that we now let  $x_1, \dots, x_k \rightarrow x_0$ . Then  $\xi \rightarrow x_0$  and

$$(41) \quad \lim_{x_1, \dots, x_k \rightarrow x_0} f[x_0, \dots, x_k] = \frac{1}{k!} f^{(k)}(x_0).$$

We can use this relation to define a divided difference where there are "coincident" nodes. For example  $f[x_0, x_1]$  when  $x_0 = x_1$  by  $f[x_0, x_0] = f'(x_0)$ , etc. This is going to be very useful for the following interpolation problem.

## 7 Hermite Interpolation

The Hermite interpolation problem is: given values of  $f$  and some of its derivatives at the nodes  $x_0, x_1, \dots, x_n$ , find the interpolating polynomial of smallest degree interpolating those values. This polynomial is called the *Hermite Interpolation Polynomial* and can be obtained with a minor modification to the Newton's form representation.

For example: Suppose we look for a polynomial of  $P$  of lowest degree which satisfies the interpolation conditions:

$$\begin{aligned} P(x_0) &= f(x_0), \\ P'(x_0) &= f'(x_0), \\ P(x_1) &= f(x_1), \\ P'(x_1) &= f'(x_1). \end{aligned}$$

We can view this problem as a limiting case of polynomial interpolation of  $f$  at two pairs of coincident nodes,  $x_0, x_0, x_1, x_1$  and we can use Newton's Interpolation form to obtain  $P$ . The table of divided differences, in view of (41), is

$$(42) \quad \begin{array}{c|ccc} x_0 & f(x_0) & & \\ x_0 & f(x_0) & f'(x_0) & \\ x_1 & f(x_1) & f[x_0, x_1] & f[x_0, x_0, x_1] \\ x_1 & f(x_1) & f'(x_1) & f[x_0, x_1, x_1] & f[x_0, x_0, x_1, x_1] \end{array}$$

and

$$P(x) = f(x_0) + f'(x_0)(x - x_0) + f[x_0, x_0, x_1](x - x_0)^2 + f[x_0, x_0, x_1, x_1](x - x_0)^2(x - x_1).$$

**Example 4.** Let  $f(0) = 1$ ,  $f'(0) = 0$  and  $f(1) = \sqrt{2}$ . Find the Hermite Interpolation Polynomial.

We construct the table of divided differences as follows:

$$(43) \quad \begin{array}{c|ccc} 0 & \textcircled{1} & & \\ 0 & 1 & \textcircled{0} & \\ 1 & \sqrt{2} & \sqrt{2} - 1 & \textcircled{\sqrt{2} - 1} \end{array}$$

and therefore

$$P(x) = 1 + 0(x - 0) + (\sqrt{2} - 1)(x - 0)^2 = 1 + (\sqrt{2} - 1)x^2.$$